

Variance in Type Systems and Variance-Based Parametric Types

Based on Igarashi and Viroli's paper from ECOOP 2002 (excellent paper! Value more in taste, than in novelty)

- This is a mechanism that got integrated in Java generics with different syntax

Subtyping

- Roughly, when a type is a subset of another
- What does that mean for method signatures? (covariance/contravariance of arguments result types)
- Consider (which one really defines a subset?):

```
interface I1 {
    Animal foo(Dog d);
}
interface I2 extends I1 {
    Dog foo(Animal d);
}
interface I3 extends I1 {
    Object foo(PrettyDog d);
}
interface I4 extends I2 {
    Dog foo(Dog d);
}
```

Variance Flavors

- Covariance: $R <: S \Rightarrow C\langle R \rangle <: C\langle S \rangle$
- Contravariance: $R <: S \Rightarrow C\langle S \rangle <: C\langle R \rangle$
- Bivariance: $C\langle R \rangle <: C\langle S \rangle$, for all R and S
- Invariance: $C\langle R \rangle <: C\langle S \rangle \Rightarrow R = S$

Question: How Can We Have Safe Variance?

Two basic principles, applied in a variety of mechanisms:

- C covariant in X means that X should not be the type of a public (and writeable—e.g., non-final) field or an argument type of a public method
- C contravariant in X means that X should not be the type of a public, readable field, or the return type of a public method

Classical, Restrictive Approach

```
class Pair<X extends Object,
        Y extends Object> {
    private X fst;
    private Y snd;
    Pair(X fst, Y snd) {
        this.fst = fst;
        this.snd = snd;
    }
    void SetFst(X fst) {
        this.fst = fst;
    }
    Y getSnd() { return snd; }
}
```

- Pair is covariant in Y , contravariant in X
- why don't constructors matter?
- E.g. $\text{Pair}\langle \text{Object}, \text{Integer} \rangle$ can be used where $\text{Pair}\langle \text{String}, \text{Number} \rangle$ is expected
- $(\text{Integer} <: \text{Number} <: \text{Object})$

Limitations of Classical Approach

Usually we use the type parameter both in covariant and in contravariant roles

```
class Vector<X> {
    private X[] ar;

    Vector(int size){ar = new X[size];}
    int size(){return ar.length;}
    X getElementAt(int i){return ar[i];}
    void setElementAt(X t,int i) {
        ar[i] = t;
    }
}
```

- Too conservative to infer variance from code

New Insight

- Instead of conservatism, disallow some uses of methods based on the statically known type information
- Think of the same single code for Vector as defining 4 classes:
 - the regular Vector
 - the covariant Vector (with only read-only methods)
 - the contravariant Vector (only write-only methods)
 - the bivariant Vector (no methods with Xs in their parameter or return list—“frozen” Vector)

Variance Annotations

Three kinds of annotations:

- + : covariance (think “const” or “read-only”)
- - : contravariance (think “write-only”)
- * : bivariance (think “contents not touched”)

Interpretation:

- $C<+T>$: the union of all invariant types of the form $C<S>$, where $S <: T$
 - C with T used only to read from
- $C<-T>$: the union of all invariant types of the form $C<S>$, where $T <: S$
 - C with T used only to write to
- $C<*>$: all invariant types of the form $C<S>$

(Note I say “union”—types are sets of values)

Rules

(For multiple type parameters, the rules apply by varying a single parameter and keeping all others the same)

$C<T> <: C<+T>$
- $\text{Vector}<\text{Integer}> <: \text{Vector}<+\text{Integer}>$

$C<T> <: C<-T>$
- $\text{Vector}<\text{Integer}> <: \text{Vector}<-\text{Integer}>$

$C<+T> <: C<*>$
- $\text{Vector}<+\text{Integer}> <: \text{Vector}<*>$

$C<-T> <: C<*>$
- $\text{Vector}<-\text{Integer}> <: \text{Vector}<*>$

$S <: T \Rightarrow C <+S> <: C<+T>$
- $\text{Vector}<+\text{Integer}> <: \text{Vector}<+\text{Number}>$

$S <: T \Rightarrow C <-T> <: C<-S>$
- $\text{Vector}<-\text{Number}> <: \text{Vector}<-\text{Integer}>$

Example Applications: Covariance

```
class Vector<X> {
    ...
    void fillFrom(Vector<+X> v, int i) {
        for (int j=i; j<v.size(); j++)
            setElementAt(
                v.getElementAt(j-i),j);
    }
}
```

Fills a vector (beginning at position *i*) by reading the contents of another vector. *v* is read-only, the method is covariant

```
Vector<Number> vn =
    new Vector<Number>(20);
Vector<Integer> vi = new
    Vector<Integer>(10);
Vector<Float> vf = new
    Vector<Float>(10);

vn.fillFrom(vi,0);
vn.fillFrom(vf,10);
```

Example Applications: Covariance

```
class Vector<X> {
    ...
    void fillFromVector(
        Vector<+Vector<+X>> vv) {
        int pos = 0;
        for (int i=0; i<vv.size(); i++) {
            Vector<+X> v = vv.getElementAt(i);
            if (pos+v.size() >= size()) break;
            fillFrom(v,pos);
            pos +=v.size();
        }
    }
}
```

Fills a vector with the contents of all vectors in a vector-of-vectors

E.g. the `Vector<X>` object (this) can be `Vector<Number>`, while `vv` is a `Vector<Vector<+Number>>` (e.g., holding a vector of Integers and a vector of floats)

Example Applications: Contravariance

```
class Vector<X> {
    ...
    void fillTo(Vector<-X> v, int i) {
        for (int j=i; j<v.size(); j++)
            v.setElementAt(
                getElementAt(j),j-i);
    }
}
```

Fills vector *v* by reading the contents of another vector (beginning at position *i*). *v* is write-only, the method is contravariant

```
Vector<Number> vn =
    new Vector<Number>(20);
Vector<Integer> vi = new
    Vector<Integer>(10);
Vector<Float> vf = new
    Vector<Float>(10);

vi.fillTo(vn,0);
vf.fillTo(vn,10);
```

Example Applications: Bivariance

```
int countVec(Vector<+Vector<*>> vv) {
    int sz = 0;
    for (int i=0; i < vv.size(); i++) {
        sz += vv.getElementAt(i).size();
    }
    return sz;
}
```

We count all elements of members of a vector-of-vectors. The second level vectors are not touched, the vector-of-vectors is only read

As another example, think of a vector of pairs, where only the first element of each pair is read and the Vector is not modified:

```
Vector<+Pair<+X, *>>
```

Assessment

- The variance annotations (which could be inferred if all the code is available for analysis) yield more generic code
- Similar to parametric (template) methods, with bounds on the template parameters
 - but need lower bounds, in addition to the usual “X extends C” (upper bound)
 - the mechanisms are complementary—each can do some things better than the other (read the paper for details!)
- Informally, parametric types with variance are like bounded existential types: e.g.,
`Vector<+C>` is like a type
`<exists X <: C> Vector<X>`